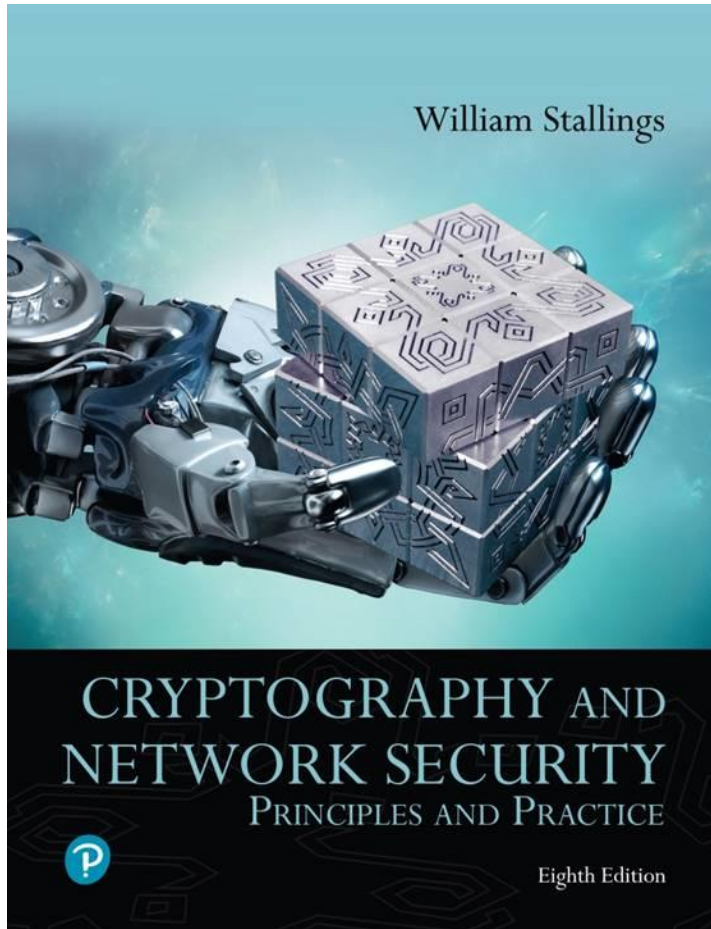


# Cryptography and Network Security: Principles and Practice

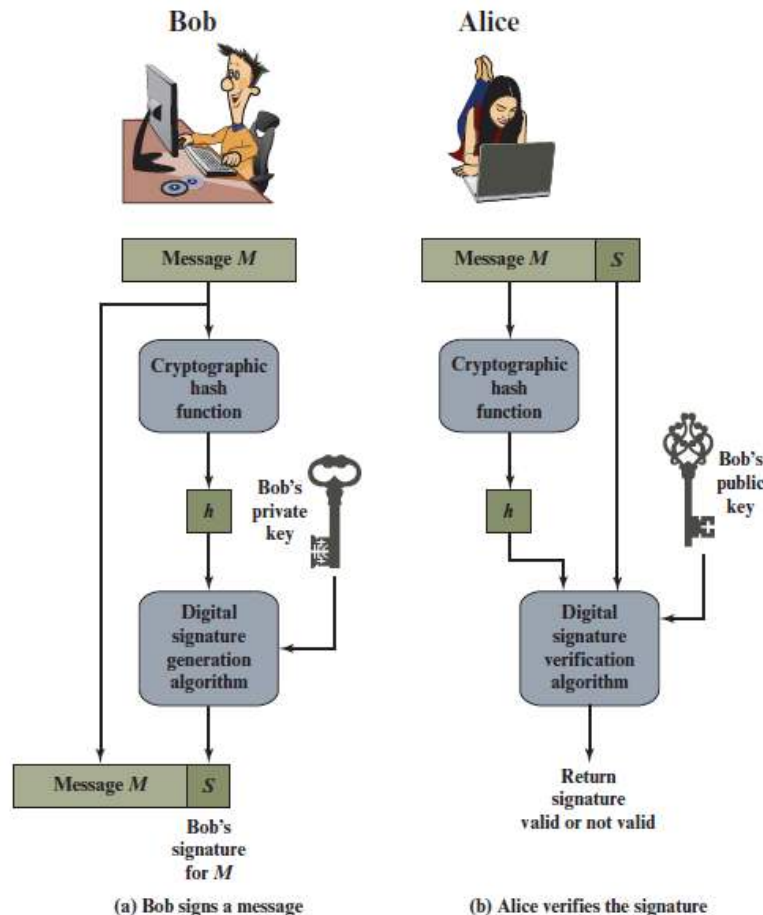
Eighth Edition



## Chapter 13

### Digital Signatures

# Figure 13.1 Simplified Depiction of Essential Elements of Digital Signature Process



# Digital Signature Properties

- It must verify the author and the date and time of the signature
- It must authenticate the contents at the time of the signature
- It must be verifiable by third parties to resolve disputes

# Attacks

- **Key-only attack**
  - C only knows A's public key
- **Known message attack**
  - C is given access to a set of messages and their signatures
- **Generic chosen message attack**
  - C chooses a list of messages before attempting to break A's signature scheme, independent of A's public key; C then obtains from A valid signatures for the chosen messages
- **Directed chosen message attack**
  - Similar to the generic attack, except that the list of messages to be signed is chosen after C knows A's public key but before any signatures are seen
- **Adaptive chosen message attack**
  - C may request from A signatures of messages that depend on previously obtained message-signature pairs

# Forgeries

- **Total break**
  - C determines A's private key
- **Universal forgery**
  - C finds an efficient signing algorithm that provides an equivalent way of constructing signatures on arbitrary messages
- **Selective forgery**
  - C forges a signature for a particular message chosen by C
- **Existential forgery**
  - C forges a signature for at least one message; C has no control over the message

# Digital Signature Requirements

- The signature must be a bit pattern that depends on the message being signed
- The signature must use some information unique to the sender to prevent both forgery and denial
- It must be relatively easy to produce the digital signature
- It must be relatively easy to recognize and verify the digital signature
- It must be computationally infeasible to forge a digital signature, either by constructing a new message for an existing digital signature or by constructing a fraudulent digital signature for a given message
- It must be practical to retain a copy of the digital signature in storage

# Direct Digital Signature

- Refers to a digital signature scheme that involves only the communicating parties
  - It is assumed that the destination knows the public key of the source
- Confidentiality can be provided by encrypting the entire message plus signature with a shared secret key
  - It is important to perform the signature function first and then an outer confidentiality function
  - In case of dispute some third party must view the message and its signature
- The validity of the scheme depends on the security of the sender's private key
  - If a sender later wishes to deny sending a particular message, the sender can claim that the private key was lost or stolen and that someone else forged his or her signature
  - One way to thwart or at least weaken this ploy is to require every signed message to include a timestamp and to require prompt reporting of compromised keys to a central authority

# ElGamal Digital Signature

- Scheme involves the use of the private key for encryption and the public key for decryption
- Global elements are a prime number  $q$  and  $a$ , which is a primitive root of  $q$
- Use private key for encryption (signing)
- Uses public key for decryption (verification)
- Each user generates their key
  - Chooses a secret key (number):  $1 < x_A < q-1$
  - Compute their public key:  $y_A = a^{x_A} \bmod q$



# Schnorr Digital Signature

- Scheme is based on discrete logarithms
- Minimizes the message-dependent amount of computation required to generate a signature
  - Multiplying a  $2n$ -bit integer with an  $n$ -bit integer
- Main work can be done during the idle time of the processor
- Based on using a prime modulus  $p$ , with  $p - 1$  having a prime factor  $q$  of appropriate size
  - Typically  $p$  is a 1024-bit number, and  $q$  is a 160-bit number

# NIST Digital Signature Algorithm

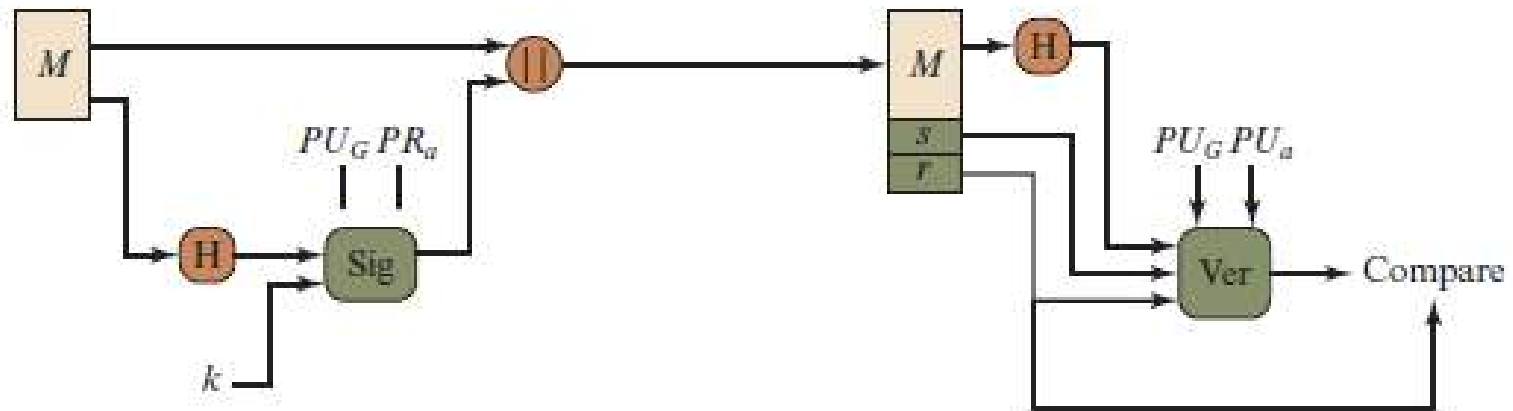
- Published by NIST as Federal Information Processing Standard FIPS 186
- Makes use of the Secure Hash Algorithm (SHA)
- The latest version, FIPS 186-3, also incorporates digital signature algorithms based on RSA and on elliptic curve cryptography



# Figure 13.2 Two Approaches to Digital Signatures



(a) RSA approach



(b) DSA approach

# Figure 13.3 The Digital Signature Algorithm (DSA)

## Global Public-Key Components

- $p$  prime number where  $2^{L-1} < p < 2^L$   
for  $512 \leq L \leq 1024$  and  $L$  a multiple of 64;  
i.e., bit length  $L$  between 512 and 1024 bits  
in increments of 64 bits
- $q$  prime divisor of  $(p - 1)$ , where  $2^{N-1} < q < 2^N$   
i.e., bit length of  $N$  bits
- $g = h(p - 1)/q$  is an exponent mod  $p$ ,  
where  $h$  is any integer with  $1 < h < (p - 1)$   
such that  $h^{(p-1)/q} \bmod p > 1$

## User's Private Key

- $x$  random or pseudorandom integer with  $0 < x < q$

## User's Public Key

$$y = g^x \bmod p$$

## User's Per-Message Secret Number

- $k$  random or pseudorandom integer with  $0 < k < q$

## Signing

$$r = (g^k \bmod p) \bmod q$$

$$s = [k^{-1} (H(M) + xr)] \bmod q$$

$$\text{Signature} = (r, s)$$

## Verifying

$$w = (s')^{-1} \bmod q$$

$$u_1 = [H(M')w] \bmod q$$

$$u_2 = (r')w \bmod q$$

$$v = [(g^{u_1} y^{u_2}) \bmod p] \bmod q$$

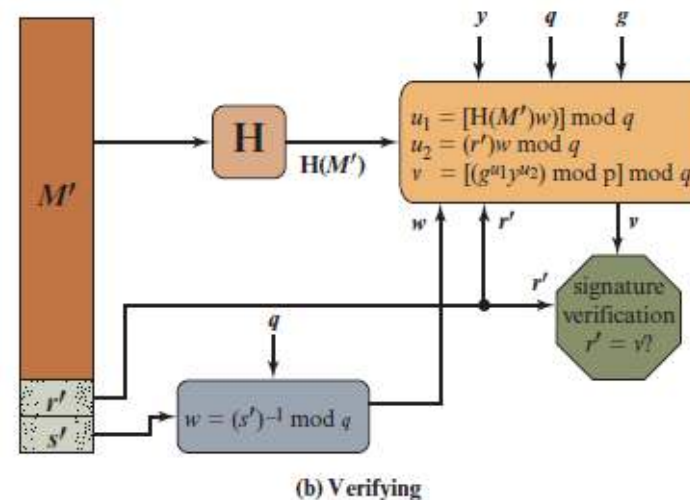
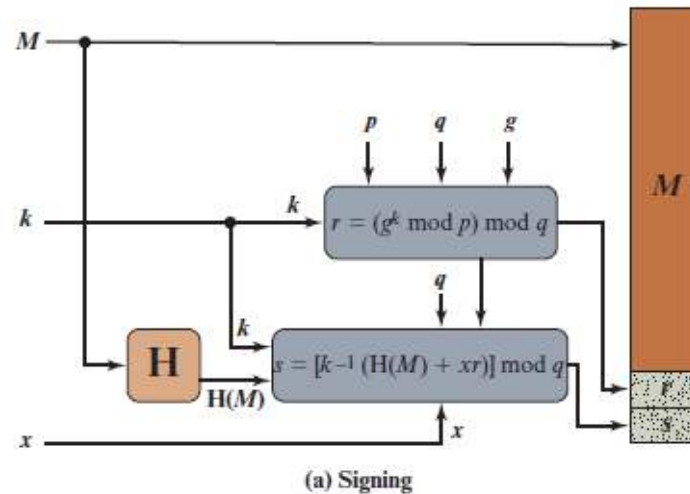
$$\text{TEST: } v = r'$$

$M$  = message to be signed

$H(M)$  = hash of  $M$  using SHA-1

$M', r', s'$  = received versions of  $M, r, s$

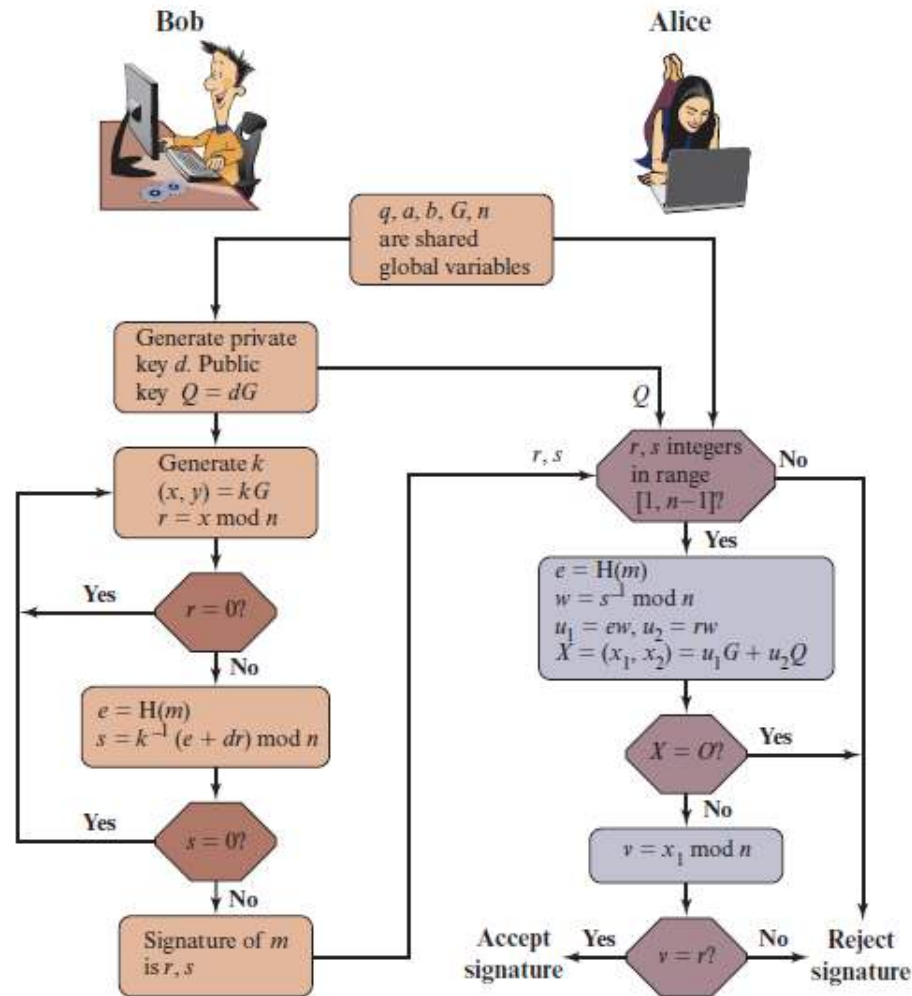
# Figure 13.4 DSA Signing and Verifying



# Elliptic Curve Digital Signature Algorithm (ECDSA)

- Four elements are involved:
  - All those participating in the digital signature scheme use the same global domain parameters, which define an elliptic curve and a point of origin on the curve
  - A signer must first generate a public, private key pair
  - A hash value is generated for the message to be signed; using the private key, the domain parameters, and the hash value, a signature is generated
  - To verify the signature, the verifier uses as input the signer's public key, the domain parameters, and the integer  $s$ ; the output is a value  $v$  that is compared to  $r$ ; the signature is verified if the  $v = r$

# Figure 13.5 ECDSA Signing and Verifying



# RSA-PSS

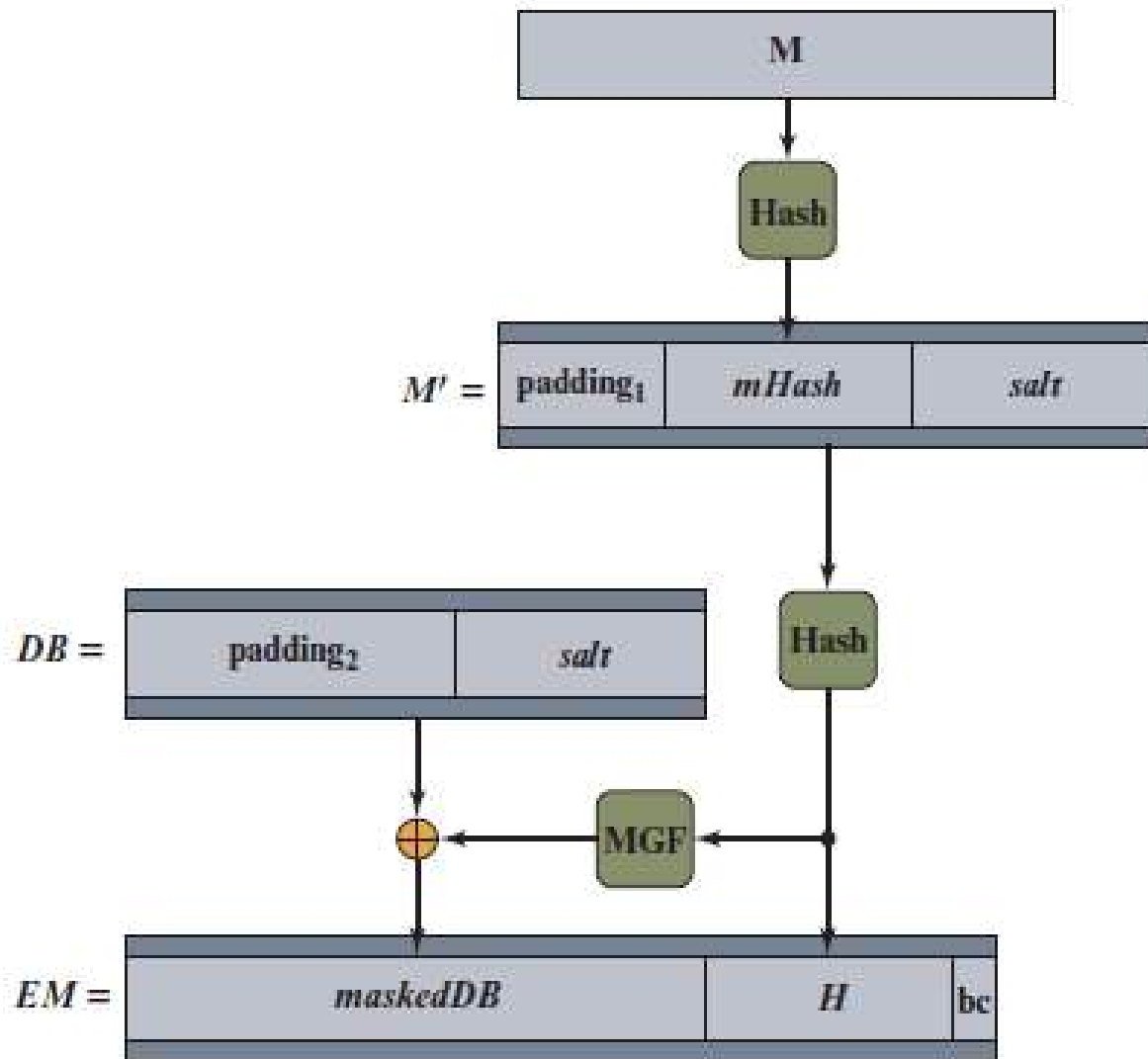
- RSA Probabilistic Signature Scheme
- Included in the 2009 version of FIPS 186
- Latest of the RSA schemes and the one that RSA Laboratories recommends as the most secure of the RSA schemes
- For all schemes developed prior to PSS it has not been possible to develop a mathematical proof that the signature scheme is as secure as the underlying RSA encryption/decryption primitive
- The PSS approach was first proposed by Bellare and Rogaway
- This approach, unlike the other RSA-based schemes, introduces a randomization process that enables the security of the method to be shown to be closely related to the security of the RSA algorithm itself



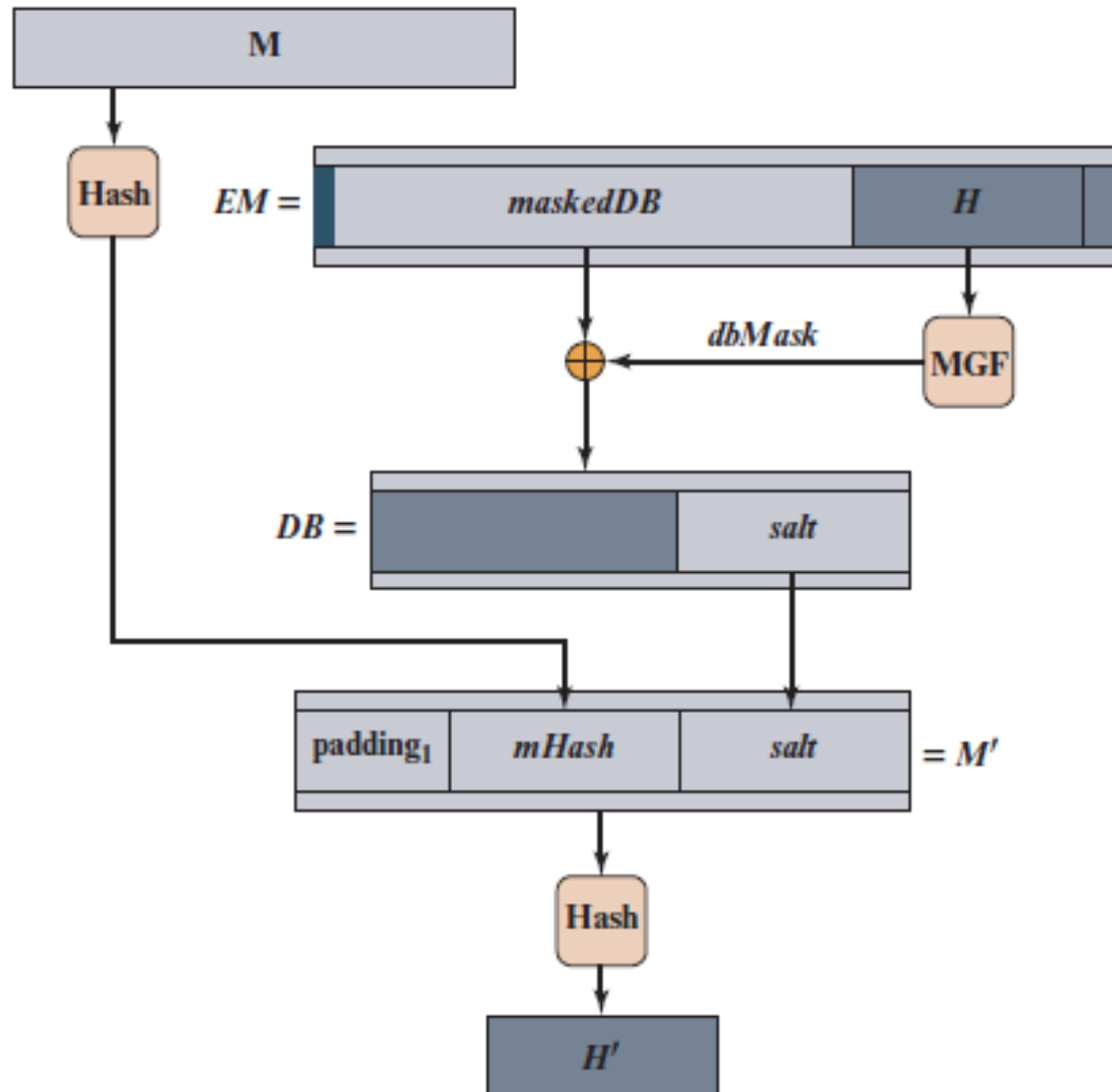
# Mask Generation Function (MGF)

- Typically based on a secure cryptographic hash function such as SHA-1
  - Is intended to be a cryptographically secure way of generating a message digest, or hash, of variable length based on an underlying cryptographic hash function that produces a fixed-length output

# Figure 13.6 RSA-PSS Encoding



# Figure 13.7 RSA-PSS EM Verification



# Summary

- Present an overview of the digital signature process
- Understand the ElGamal digital signature scheme
- Understand the Schnorr digital signature scheme
- Understand the NIST digital signature scheme
- Compare and contrast the NIST digital signature scheme with the ElGamal and Schnorr digital signature schemes
- Understand the elliptic curve digital signature scheme
- Understand the RSA-PSS digital signature scheme



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